CSCI 5451: Introduction to Parallel Computing

Lecture 25: Histograms in Cuda



Announcements (12/03)

- Updated HW4
 - Question 4 in HW4 has been updated as access to Colab is restricted for some
 - Be sure to review this updated file
 - The question is largely the same, but the timings are now given explicitly (you do not need to run separately on an A100 + T4)
- Cuda Machines now fully available except for GPU 2
 - o If you are assigned to GPU 2, then select one of the other gpus to use at random



Overview

- Histograms
- Parallel Histograms
- Simple histogram Kernel
- Anatomy of Atomic Updating in Cuda
- Improved Histogram Kernels (Privatization, Shared Memory, Coarsening, Aggregation)

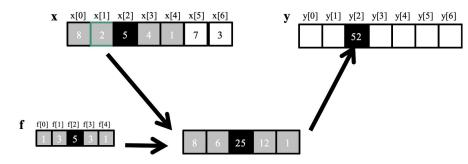


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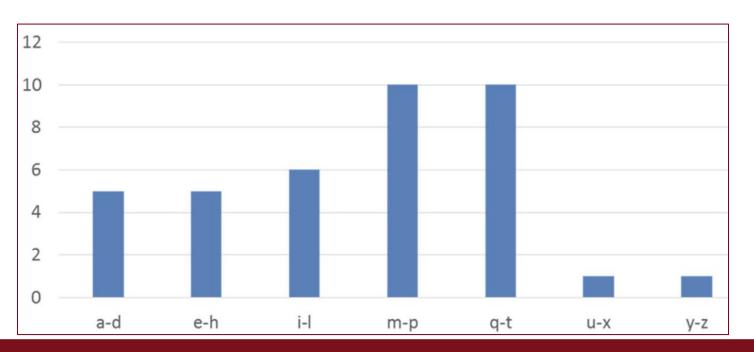


- So far, the algorithms we have examined have used an output decomposition
- ☐ Each thread writes to a separate location (owner-computes)
- There is no contention for writing to the same memory location between different threads
- We do not need to coordinate writes across threads at all - they are independent





We will use histograms to demonstrate a case where multiple threads need to write to the same location at a given time





- ☐ The below algorithm is serial (CPU), runs in O(N) and assumes *data* is in ACSII format
- ☐ The result stores the number of groups of 4 chars ('a-d', 'e-h', 'i-l', etc.) in *histo*
- This takes reasonable advantage of the cache by operating on consecutive entries of data



How to parallelize?

- ☐ The below algorithm is serial (CPU), runs in O(N) and assumes *data* is in ACSII format
- ☐ The result stores the number of groups of 4 chars ('a-d', 'e-h', 'i-l', etc.) in *histo*
- This takes reasonable advantage of the cache by operating on consecutive entries of data

How to parallelize? Input Decomposition

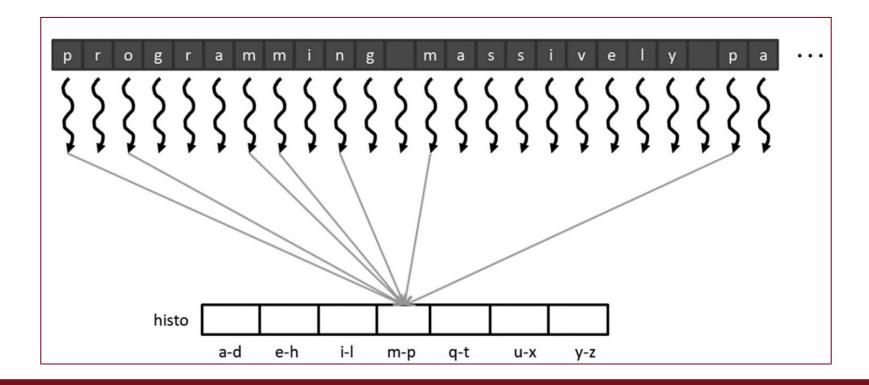
- ☐ The below algorithm is serial (CPU), runs in O(N) and assumes *data* is in ACSII format
- ☐ The result stores the number of groups of 4 chars ('a-d', 'e-h', 'i-l', etc.) in *histo*
- This takes reasonable advantage of the cache by operating on consecutive entries of data

Overview

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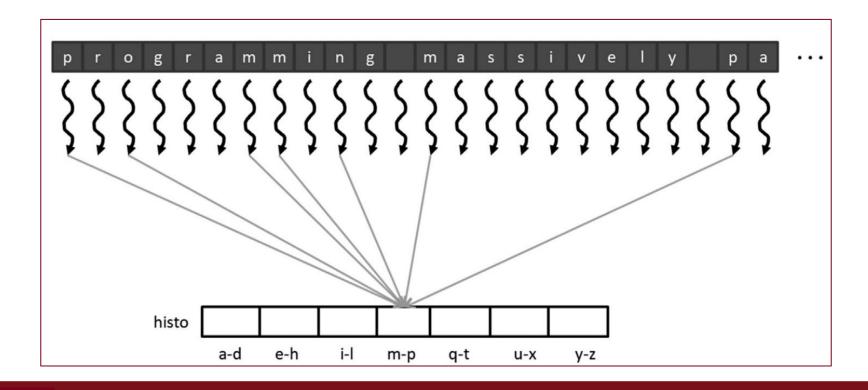
Parallel Histogram





Parallel Histogram

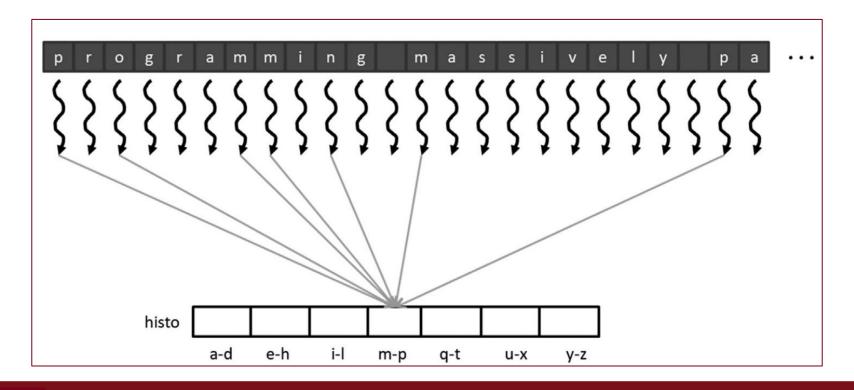
What problem does this introduce?





Parallel Histogram

What problem does this introduce? Race conditions





Race Conditions

- Recall that race conditions can occur in our program when we are issuing concurrent writes to the same memory location
- The exact ordering of instructions may result in our program not updating variables as a serial version of our program initially would have

Time	Thread 1	Thread 2
1	(0) Old \leftarrow histo[x]	
2	(1) New ← Old + 1	
3		(0) Old \leftarrow histo[x]
4	(1) histo[x] \leftarrow New	
5		(1) New ← Old + 1
6		(1) histo[x] \leftarrow New



Overview

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Atomic Additions

As is done within OpenMP and pthreads, we can make use of an atomic addition to ensure that our writes will occur sequentially when multiple threads attempt to write to the same location (this eliminates the race condition)

int atomicAdd(int* address, int val);



Histogram Kernel

```
01
      global void histo kernel (char *data, unsigned int length,
                                  unsigned int *histo) {
02
        unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
03
        if (i < length) {
04
            int alphabet position = data[i] - 'a';
05
            if (alphabet position >= 0 && alpha position < 26) {
06
                atomicAdd(&(histo[alphabet position/4]), 1);
07
08
09
```



Histogram Kernel

How should we launch this kernel? Are we coalescing?



Overview

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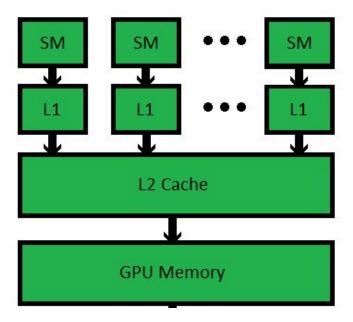


How do atomic additions work in practice?

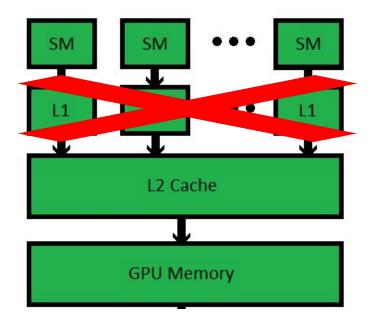


- All threads within a warp execute the same instruction for the atomicAdd
- ☐ Unlike non-atomic memory transactions, atomic memory transactions skip the L1 cache altogether and go straight to the L2 global cache. Why?

0



- All threads within a warp execute the same instruction for the atomicAdd
- ☐ Unlike non-atomic memory transactions, atomic memory transactions skip the L1 cache altogether and go straight to the L2 global cache. Why?
 - Global variables are shared. We need the transactions to them to be sequential across all threads in the grid, not just a given threadblock





- Once at the L2 cache, the actual updates take place for the threads within that warp
- ☐ If multiple threads within a warp execute an update to the same location, they are *serialized* by the L2 memory controller
- □ These serialized updates are pipelined to ensure there are still some speedups

// 32 threads in a warp all execute:
atomicAdd(&shared_counter, 1);

```
Warp Pipeline after reaching L2 controller
      [issue]
      [read]
      [add]
      [write]
      [complete]
Pipeline of execution for this warp after reaching
L2 memory controller (given that all writes are to
the same ocation → shared counter)
   Thread 1: [-pipeline-]
   Thread 0: [-pipeline-]
   Thread 2:
                 [-pipeline-]
// The operations overlap in the pipeline
```

- □ The warp only continues execution after the final thread finishes executing its pipeline (and the L2 memory controller sends a signal the warp has completed its updates)
- ☐ The ordering of sequential updates is unknown

```
// All 32 threads in a warp execute the following
// Assume only one 32 thread block is dispatched
atomicAdd(&counter, threadIdx.x);
// Possible execution orders:
// Run 1: threads execute in order 0,1,2,3,...,31
// Run 2: threads execute in order 5,2,31,0,17,...
// Run 3: threads execute in some other order
// All are valid
// Final result is always: counter += (0+1+2+...+31) = 496
// But intermediate states are unpredictable
```



What about threads in the same threadblock, but different warps? What about between different threadblocks?



- As within warps, transactions to the same memory locations will always be serial in threads from different warps, and the ordering remains unknown
- In total, we cannot know which threads in a warp, warps in a block, or blocks in a grid will execute their transactions first
- The exact order of atomic additions by individual thread is unknown

```
atomicAdd(&global_counter, 1);
```

```
Pipeline stage for the entire warp:
       [Issue] \rightarrow
       [L2 request] \rightarrow
       [L2 queue] \rightarrow
       [read] \rightarrow
       [modify] \rightarrow
       [write] \rightarrow
       [respond]
Multiple atomics of different warps can be
in flight concurrently:
       Warp 1, Block 0: [-pipeline-]
       Warp 1, Block 2: [—pipeline—]
       Warp 0, Block 0:
                                 [-pipeline-]
```



- ☐ Different warps *do not implicitly synchronize* their updates
- Likewise, warps across different threadblocks do not synchronize their updates
- We can synchronize the value within a threadblock with __synchthreads()
- There is no way to synchronize the value between threadblocks at a given point during kernel execution

```
// BAD - race condition within block
atomicAdd(&counter, 1);
if (counter > 50) { ... } // Undefined behavior!

// GOOD - synchronized within block
// - Not synchronized between blocks!!
atomicAdd(&counter, 1);
__syncthreads();
if (counter > 50) { ... } // Now safe within block
```



Utilizing this Anatomy in Practice

In general, atomic operations are useful, provided your kernel has the following properties

- The variable being atomically updated is write-only, and the operations performed on this variable are independent of the ordering
- □ **Or**, if you are reading the atomic variable during program execution:
 - Threads in different threadblocks are accessing different regions
 - And you are synchronizing after issuing your writes to ensure coherence within each threadblock



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Revisiting the Original Kernel

Why is this kernel not very fast?

```
01
     global void histo kernel (char *data, unsigned int length,
                                  unsigned int *histo)
02
        unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
03
        if (i < length)
04
            int alphabet position = data[i] - 'a';
05
            if (alphabet position >= 0 && alpha position < 26) {
06
                atomicAdd(&(histo[alphabet position/4]), 1);
07
08
09
```



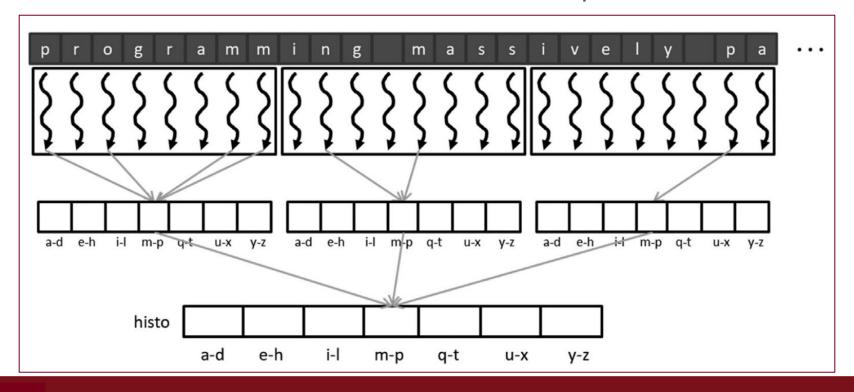
Revisiting the Original Kernel

Why is this kernel not very fast? Excessive serialization leads to idle warps

```
01
      global void histo kernel (char *data, unsigned int length,
                                  unsigned int *histo)
02
        unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
03
        if (i < length)
04
            int alphabet position = data[i] - 'a';
05
            if (alphabet position >= 0 && alpha position < 26) {
                atomicAdd(&(histo[alphabet position/4]), 1);
06
07
08
09
```



Copy the *histo* array for each block of threads that we will spawn





Copy the *histo* array for each block of threads that we will spawn

```
01
      global void histo private kernel (char *data, unsigned int length,
                                           unsigned int *histo) {
      unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
02
03
      if(i < length) {
04
        int alphabet position = data[i] - 'a';
05
        if (alphabet position >= 0 && alphabet position < 26) {
06
           atomicAdd(&(histo[blockIdx.x*NUM BINS + alphabet position/4]), 1);
07
08
09
      if(blockIdx.x > 0) {
10
          syncthreads();
        for(unsigned int bin=threadIdx.x; bin<NUM BINS; bin += blockDim.x) {</pre>
11
12
           unsigned int binValue = histo[blockIdx.x*NUM BINS + bin];
13
           if(binValue > 0) {
14
              atomicAdd(&(histo[bin]), binValue);
15
16
17
18
```



Does this use memory coalescing? Why is this advantageous?

```
01
      global void histo private kernel (char *data, unsigned int length,
                                          unsigned int *histo) {
      unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
02
03
      if(i < length) {
04
        int alphabet position = data[i] - 'a';
05
        if (alphabet position >= 0 && alphabet position < 26) {
06
           atomicAdd(&(histo[blockIdx.x*NUM BINS + alphabet position/4]), 1);
07
08
09
      if(blockIdx.x > 0) {
10
          syncthreads();
11
        for (unsigned int bin=threadIdx.x; bin<NUM BINS; bin += blockDim.x) {
12
           unsigned int binValue = histo[blockIdx.x*NUM BINS + bin];
13
           if(binValue > 0) {
14
              atomicAdd(&(histo[bin]), binValue);
15
16
17
18
```



Do all threads participate in updating the full histogram? Where is this full histogram stored?

```
01
      global void histo private kernel (char *data, unsigned int length,
                                           unsigned int *histo) {
      unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
02
03
      if(i < length) {
04
        int alphabet position = data[i] - 'a';
05
        if (alphabet position >= 0 && alphabet position < 26) {
06
           atomicAdd(&(histo[blockIdx.x*NUM BINS + alphabet position/4]), 1);
07
08
09
      if(blockIdx.x > 0) {
10
          syncthreads();
        for(unsigned int bin=threadIdx.x; bin<NUM BINS; bin += blockDim.x) {</pre>
11
12
           unsigned int binValue = histo[blockIdx.x*NUM BINS + bin];
13
           if(binValue > 0) {
14
              atomicAdd(&(histo[bin]), binValue);
15
16
17
18
```



```
global
              void histo private kernel (char* data, unsigned int length,
                                         unsigned int* histo) {
02
         // Initialize privatized bins
03
                unsigned int histo s[NUM BINS];
04
     for (unsigned int bin = threadIdx.x; bin< NUM BINS; bin += blockDim.x)
05
        histo s[bin] = 0u;
06
07
       syncthreads();
08
        // Histogram
09
     unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
10
      if(i < length) {
        int alphabet position = data[i] - 'a';
11
        if (alphabet position >= 0 && alphabet position < 26) {
12
13
           atomicAdd(&(histo s[alphabet position/4]), 1);
14
15
16
        syncthreads();
17
        // Commit to global memory
18
      for (unsigned int bin=threadIdx.x; bin<NUM BINS; bin += blockDim.x) {
19
            unsigned int binValue = histo s[bin];
20
            if(binValue > 0) {
21
              atomicAdd(&(histo[bin]), binValue);
22
23
24
```



Do all threads update the shared memory values?

```
global
              void histo private kernel (char* data, unsigned int length,
                                          unsigned int* histo) {
02
         // Initialize privatized bins
03
                unsigned int histo s[NUM BINS];
04
     for (unsigned int bin = threadIdx.x; bin< NUM BINS; bin += blockDim.x)
05
        histo s[bin] = 0u;
06
07
       syncthreads();
08
        // Histogram
09
     unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
      if(i < length) {
10
        int alphabet position = data[i] - 'a';
        if(alphabet position >= 0 && alphabet position < 26) {
13
           atomicAdd(&(histo s[alphabet position/4]), 1);
14
15
16
        syncthreads();
17
        // Commit to global memory
18
      for(unsigned int bin=threadIdx.x; bin<NUM BINS; bin += blockDim.x) {</pre>
19
            unsigned int binValue = histo s[bin];
20
            if(binValue > 0) {
21
              atomicAdd(&(histo[bin]), binValue);
22
23
24
```



Why is updating the shared memory locations better?

```
global
              void histo private kernel (char* data, unsigned int length,
                                          unsigned int* histo) {
02
         // Initialize privatized bins
03
                unsigned int histo s[NUM BINS];
04
     for (unsigned int bin = threadIdx.x; bin< NUM BINS; bin += blockDim.x)
05
        histo s[bin] = 0u;
06
07
       syncthreads();
08
        // Histogram
09
     unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
10
      if(i < length) {
        int alphabet position = data[i] - 'a';
        if(alphabet position >= 0 && alphabet position < 26) {
13
           atomicAdd(&(histo s[alphabet position/4]), 1);
14
15
16
        syncthreads();
17
        // Commit to global memory
18
      for(unsigned int bin=threadIdx.x; bin<NUM BINS; bin += blockDim.x) {</pre>
19
            unsigned int binValue = histo s[bin];
20
            if(binValue > 0) {
21
              atomicAdd(&(histo[bin]), binValue);
22
23
24
```



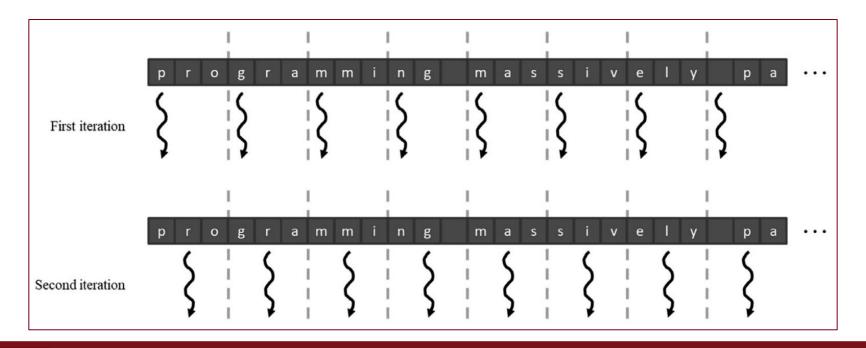
Do all threads write to global memory?

```
global
              void histo private kernel (char* data, unsigned int length,
                                          unsigned int* histo) {
02
         // Initialize privatized bins
03
                unsigned int histo s[NUM BINS];
04
     for (unsigned int bin = threadIdx.x; bin< NUM BINS; bin += blockDim.x)
05
        histo s[bin] = 0u;
06
07
       syncthreads();
08
        // Histogram
09
     unsigned int i = blockIdx.x*blockDim.x + threadIdx.x;
10
      if(i < length) {
        int alphabet position = data[i] - 'a';
11
12
        if(alphabet position >= 0 && alphabet position < 26) {
13
           atomicAdd(&(histo s[alphabet position/4]), 1);
14
15
16
        syncthreads();
17
        // Commit to global memory
18
      for (unsigned int bin=threadIdx.x; bin<NUM BINS; bin += blockDim.x) {
19
            unsigned int binValue = histo s[bin];
20
            if(binValue > 0) {
21
              atomicAdd(&(histo[bin]), binValue);
22
23
24
```



Coarsening Kernel (No 1)

We can add coarsening by giving each thread a set of input elements to work on which are sequential





Coarsening Kernel (No 1)

```
global void histo private kernel (char* data, unsigned int length,
01
                                          unsigned int* histo) {
02
        // Initialize privatized bins
03
        shared unsigned int histo s[NUM BINS];
      for(unsigned int bin = threadIdx.x; bin<NUM BINS; bin += blockDim.x)</pre>
04
05
          histo s[binIdx] = 0u;
06
07
         syncthreads();
08
        // Histogram
09
      unsigned int tid = blockIdx.x*blockDim.x + threadIdx.x;
10
      for (unsigned int i=tid*CFACTOR; i<min((tid+1)*CFACTOR, length); ++i) {
11
          int alphabet position = data[i] - 'a';
          if(alphabet position >= 0 && alphabet position < 26) {
12
13
                atomicAdd(&(histo s[alphabet position/4]), 1);
14
15
16
        syncthreads();
17
        // Commit to global memory
18
      for(unsigned int bin = threadIdx.x; bin<NUM BINS; bin += blockDim.x) {</pre>
19
          unsigned int binValue = histo s[binIdx];
20
          if(binValue > 0) {
21
            atomicAdd(&(histo[binIdx]), binValue);
22
23
24
```



Coarsening Kernel (No 1)

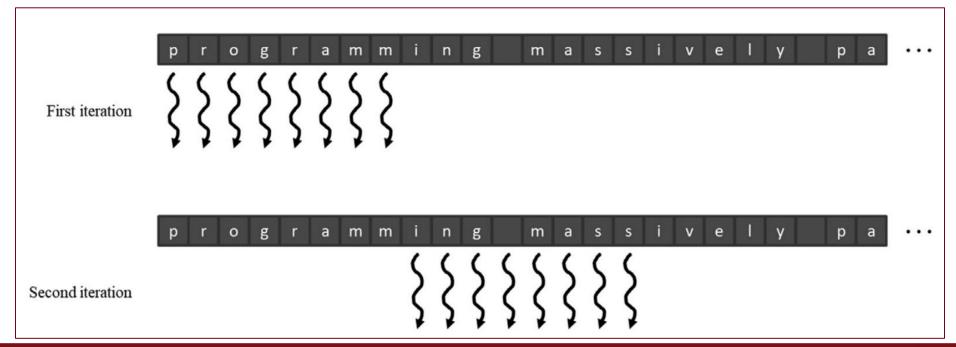
Why is this a poor coarsening implementation?

```
01
      global void histo private kernel (char* data, unsigned int length,
                                           unsigned int* histo) {
02
        // Initialize privatized bins
03
               unsigned int histo s[NUM BINS];
      for(unsigned int bin = threadIdx.x; bin<NUM BINS; bin += blockDim.x)</pre>
04
05
          histo s[binIdx] = 0u;
06
07
         syncthreads();
08
        // Histogram
09
      unsigned int tid = blockIdx.x*blockDim.x + threadIdx.x;
10
      for (unsigned int i=tid*CFACTOR; i<min((tid+1)*CFACTOR, length); ++i) {
11
          int alphabet position = data[i] - 'a';
          if (alphabet position >= 0 && alphabet position < 26) {
13
                atomicAdd(&(histo s[alphabet position/4]), 1);
14
15
16
        syncthreads();
17
        // Commit to global memory
18
      for(unsigned int bin = threadIdx.x; bin<NUM BINS; bin += blockDim.x)</pre>
19
          unsigned int binValue = histo s[binIdx];
20
          if(binValue > 0) {
            atomicAdd(&(histo[binIdx]), binValue);
22
23
24
```



Coarsening Kernel (No 2)

We now add a coalesced version of coarsening that allows for more efficient memory loading.





Coarsening Kernel (No 2)

```
global void histo private kernel (char* data, unsigned int length,
01
                                          unsigned int* histo) {
02
        // Initialize privatized bins
03
        shared unsigned int histo s[NUM BINS];
04
      for (unsigned int bin=threadIdx.x; bin<NUM BINS; bin += blockDim.x) {
05
            histo s[binIdx] = 0u;
06
07
        syncthreads();
08
        // Histogram
09
      unsigned int tid = blockIdx.x*blockDim.x + threadIdx.x;
10
      for(unsigned int i = tid; i < length; i += blockDim.x*gridDim.x) {</pre>
11
         int alphabet position = data[i] - 'a';
12
         if(alphabet position >= 0 && alphabet position < 26) {
13
             atomicAdd(&(histo s[alphabet position/4]), 1);
14
15
16
        syncthreads();
17
        // Commit to global memory
18
      for(unsigned int bin = threadIdx.x; bin<NUM BINS; bin += blockDim.x)</pre>
19
         unsigned int binValue = histo s[binIdx];
20
         if(binValue > 0) {
21
            atomicAdd(&(histo[binIdx]), binValue);
22
23
24
```



Aggregation Kernel

```
global void histo private kernel (char* data, unsigned int length,
                                          unsigned int* histo) {
        // Initialize privatized bins
03
        shared unsigned int histo s[NUM BINS];
04
      for(unsigned int bin = threadIdx.x; bin < NUM BINS; bin += blockDim.x){</pre>
05
          histo s[bin] = 0u;
06
07
        syncthreads();
        // Histogram
09
      unsigned int accumulator = 0;
10
      int prevBinIdx = -1;
11
      unsigned int tid = blockIdx.x*blockDim.x + threadIdx.x;
      for(unsigned int i = tid; i < length; i += blockDim.x*gridDim.x) {</pre>
13
         int alphabet position = data[i] - 'a';
14
         if(alphabet position >= 0 && alphabet position < 26) {
15
            int bin = alphabet position/4;
16
            if(bin == prevBinIdx) {
                 ++accumulator;
            } else {
19
                 if(accumulator > 0) {
20
                     atomicAdd(&(histo s[prevBinIdx]), accumulator);
21
                 accumulator = 1;
23
                 prevBinIdx = bin;
24
25
26
       if(accumulator > 0) {
28
            atomicAdd(&(histo s[prevBinIdx]), accumulator);
29
30
         syncthreads();
31
       // Commit to global memory
32
       for(unsigned int bin = threadIdx.x; bin<NUM BINS; bin += blockDim.x) {
            unsigned int binValue = histo s[bin];
34
            if(binValue > 0) {
35
              atomicAdd(&(histo[bin]), binValue);
36
37
```



Aggregation Kernel

When is this kernel more useful?

```
global void histo private kernel (char* data, unsigned int length,
                                          unsigned int* histo) {
        // Initialize privatized bins
03
        shared unsigned int histo s[NUM BINS];
      for(unsigned int bin = threadIdx.x; bin < NUM BINS; bin += blockDim.x) {</pre>
05
          histo s[bin] = 0u;
06
07
        syncthreads();
       // Histogram
09
      unsigned int accumulator = 0;
10
      int prevBinIdx = -1;
      unsigned int tid = blockIdx.x*blockDim.x + threadIdx.x;
      for(unsigned int i = tid; i < length; i += blockDim.x*gridDim.x) {
         int alphabet position = data[i] - 'a';
         if(alphabet position >= 0 && alphabet position < 26) {
            int bin = alphabet position/4;
            if(bin == prevBinIdx) {
                 ++accumulator;
            } else {
19
                 if(accumulator > 0) {
                    atomicAdd(&(histo s[prevBinIdx]), accumulator);
                 accumulator = 1;
                 prevBinIdx = bin;
24
25
26
       if(accumulator > 0) {
            atomicAdd(&(histo s[prevBinIdx]), accumulator);
29
30
         syncthreads();
31
       // Commit to global memory
       for(unsigned int bin = threadIdx.x; bin<NUM BINS; bin += blockDim.x) {
            unsigned int binValue = histo s[bin];
34
            if(binValue > 0) {
35
              atomicAdd(&(histo[bin]), binValue);
36
37
```



Aggregation Kernel

When is this kernel more useful?
When there are a large number of entries in data which are the same.

```
global void histo private kernel (char* data, unsigned int length,
                                          unsigned int* histo) {
        // Initialize privatized bins
03
        shared unsigned int histo s[NUM BINS];
      for(unsigned int bin = threadIdx.x; bin < NUM BINS; bin += blockDim.x) {</pre>
05
          histo s[bin] = 0u;
06
        syncthreads();
       // Histogram
      unsigned int accumulator = 0;
      int prevBinIdx = -1;
      unsigned int tid = blockIdx.x*blockDim.x + threadIdx.x;
      for(unsigned int i = tid; i < length; i += blockDim.x*gridDim.x) {
         int alphabet position = data[i] - 'a';
         if(alphabet position >= 0 && alphabet position < 26) {
            int bin = alphabet position/4;
            if(bin == prevBinIdx) {
                 ++accumulator;
            } else {
                 if(accumulator > 0) {
                    atomicAdd(&(histo s[prevBinIdx]), accumulator);
                 accumulator = 1;
                 prevBinIdx = bin;
       if(accumulator > 0) {
            atomicAdd(&(histo s[prevBinIdx]), accumulator);
29
30
         syncthreads();
31
       // Commit to global memory
       for (unsigned int bin = threadIdx.x; bin < NUM BINS; bin += blockDim.x)
            unsigned int binValue = histo s[bin];
            if(binValue > 0) {
              atomicAdd(&(histo[bin]), binValue);
36
```

